

Cone Representations, Languages, and Compilers for Convex Optimization

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Outline

- Convex optimization
- Constructive convex analysis
- Cone programming and representations
- Transforming to cone program
- Parser/solvers
- Code generation

Convex optimization problem — standard form

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & Ax = b \end{array}$$

with variable $x \in \mathbf{R}^n$

- objective and inequality constraint functions f_0, \dots, f_m are convex
- equality constraints are linear
- examples:
 - least-squares, least-squares with ℓ_1 regularization
 - linear program (LP), quadratic program (QP)
 - maximum entropy and related problems

Why convex optimization?

- beautiful, fairly complete, and useful theory
 - solution algorithms that work well in theory and practice
 - **many applications** recently discovered in
 - control
 - combinatorial optimization
 - signal and image processing
 - communications, networks
 - circuit design
 - machine learning, statistics
 - finance
- . . . and many more

How do you solve a convex problem?

- use someone else's ('standard') solver (LP, QP, SDP, . . .)
 - easy, but your problem *must* be in a standard form
 - cost of solver development amortized across many users
- write your own (custom) solver
 - lots of work, but can take advantage of special structure
- transform your problem into a standard form, and use a standard solver
 - extends reach of problems that can be solved using standard solvers
 - transformation can be hard to find, cumbersome to carry out

this talk: methods to formalize and automate the last approach

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How can you tell if a problem is convex?

need to check convexity of a function

approaches:

- use basic definition, first or second order conditions, *e.g.*, $\nabla^2 f(x) \succeq 0$
- via convex calculus: construct f using
 - library of basic functions that are convex
 - calculus rules or transformations that preserve convexity

Convex functions: Basic examples

- x^p ($p \geq 1$ or $p \leq 0$), $-x^p$ ($0 \leq p \leq 1$)
- e^x , $-\log x$, $x \log x$
- $a^T x + b$
- $x^T P x$ ($P \succeq 0$)
- $\|x\|$ (any norm)
- $\max(x_1, \dots, x_n)$

Convex functions: Less basic examples

- $x^T x / y$ ($y > 0$), $x^T Y^{-1} x$ ($Y \succ 0$)
- $\log(e^{x_1} + \dots + e^{x_n})$
- $-\log \Phi(x)$ (Φ is Gaussian CDF)
- $\log \det X^{-1}$ ($X \succ 0$)
- $\lambda_{\max}(X)$ ($X = X^T$)

Calculus rules

- *nonnegative scaling*: f convex, $\alpha \geq 0 \implies \alpha f$ convex
- *sum*: f, g convex $\implies f + g$ convex
- *affine composition*: f convex $\implies f(Ax + b)$ convex
- *pointwise maximum*: f_1, \dots, f_m convex $\implies \max_i f_i(x)$ convex
- *partial minimization*: $f(x, y)$ convex $\implies \inf_y f(x, y)$ convex
- *composition*: h convex increasing, f convex $\implies h(f(x))$ convex
- *perspective transformation*: f convex $\implies tf(x/t)$ convex for $t > 0$

Examples

from basic functions and calculus rules, we can show convexity of . . .

- piecewise-linear function: $\max_{i=1,\dots,k} (a_i^T x + b_i)$
- ℓ_1 -regularized least-squares cost: $\|Ax - b\|_2^2 + \lambda \|x\|_1$, with $\lambda \geq 0$
- sum of largest k elements of x : $x_{[1]} + \dots + x_{[k]}$
- distance to convex set C : $\mathbf{dist}(x, C) = \inf_{y \in C} \|x - y\|_2$

A general composition rule

- $h(f_1(x), \dots, f_k(x))$ is convex if h is, and for each i ,
 - f_i is affine, or
 - f_i is convex and h is nondecreasing in its i th arg, or
 - f_i is concave and h is nonincreasing in its i th arg
- this one rule subsumes most of the others
- in turn, it can be derived from the partial minimization rule

Constructive convexity verification

- build parse tree for function (expression)
- leaves are variables or constants
- nodes are composition functions of children, following general rule
- example: $(x - y)^2 / (1 - \max(x, y))$ is convex (for $x < 1, y < 1$)
 - (leaves) $x, y,$ and 1 are affine functions
 - $\max(x, y)$ is convex; $x - y$ is affine
 - $1 - \max(x, y)$ is concave
 - function u^2/v is convex, monotone decreasing in v for $v > 0$
hence, get convex function with $u = x - y, v = 1 - \max(x, y)$

Disciplined convex program

- convex optimization problem described as
 - objective: minimize (cvx expr) or maximize (ccv expr)
 - inequality constraints: cvx expr \leq ccv expr or ccv expr \geq ccv expr
 - equality constraints: aff expr = aff expr
- (convex, concave, affine) expressions formed from constants, variables, and functions using general composition rule
- functions come from a library, with known convexity, monotonicity properties
- DCP is convex-by-construction (cf. posterior convexity analysis)

(Automatic) parsing of DCP

- it's (relatively) easy to parse a DCP, given function library
- DCP is 'syntactically convex'; convexity hinges only on convexity, monotonicity attributes of functions, not their detailed meaning
- gives basic method for problem convexity detection/certification
- we'll see later another use of the resulting parse trees . . .

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Convex optimization problem — conic form

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \in \mathcal{K} \end{array}$$

with variable $x \in \mathbf{R}^n$

- objective is linear
- constraints are linear equalities and (generalized) nonnegativity
- \mathcal{K} is convex cone
- examples:
 - LP: $\mathcal{K} = \mathbf{R}_+^n$
 - semidefinite program (SDP): $\mathcal{K} = \mathbf{S}_+^n$ (PSD matrices)

Cone programming

- symmetric cone programming: \mathcal{K} is product of
 - \mathbf{R}^k ('unconstrained variables')
 - nonnegative orthant \mathbf{R}_+^k
 - Lorentz cones $\mathcal{L}^k = \{(z, t) \in \mathbf{R}^k \times \mathbf{R} \mid \|z\|_2 \leq t\}$
 - semidefinite cones \mathbf{S}_+^k

of various dimensions

- exponential cone: $\mathcal{E} = \{(x, y, t) \mid \exp(x/t) \leq y/t, t > 0\}$
- with these cones, can express almost any convex problem that arises in applications (we'll see how, shortly)

Cone programming solvers

- theory, algorithms for cone programming well developed in last 10 years
- software for symmetric cone programming widely available and used
 - SeDuMi, SDPT3 (open source; Matlab/C)
 - CSDP, SDPA (open source; C)
 - CVXOPT (open source; Python/C)
 - MOSEK, PENSDP (commercial)
- our goal: solve convex optimization problems by reduction to cone programs

Cone representation

(Nesterov, Nemirovsky)

cone representation of (convex) function f :

- $f(x)$ is optimal value of cone program

$$\begin{array}{ll} \text{minimize} & c^T x + d^T y + e \\ \text{subject to} & A \begin{bmatrix} x \\ y \end{bmatrix} = b, \quad \begin{bmatrix} x \\ y \end{bmatrix} \in \mathcal{K} \end{array}$$

- cone program in (x, y) , we but minimize only over y
- *i.e.*, we define f by optimizing over *some* variables in a cone program

Examples

- $f(x) = -(xy)^{1/2}$ is optimal value of SDP

$$\begin{array}{ll} \text{minimize} & -t \\ \text{subject to} & \begin{bmatrix} x & t \\ t & y \end{bmatrix} \succeq 0 \end{array}$$

with variable t

- $f(x, y) = x^T x / y$ is optimal value of SDP

$$\begin{array}{ll} \text{minimize} & t \\ \text{subject to} & \begin{bmatrix} tI & x \\ x^T & y \end{bmatrix} \succeq 0 \end{array}$$

with variable t

- $f(x) = x_{[1]} + \cdots + x_{[k]}$ is optimal value of LP

$$\begin{array}{ll} \text{minimize} & \mathbf{1}^T \lambda - k\nu \\ \text{subject to} & x + \nu \mathbf{1} = \lambda - \mu \\ & \lambda \succeq 0, \quad \mu \succeq 0 \end{array}$$

with variables λ, μ, ν

- $f(p, q) = p \log(p/q)$ is optimal value of exponential cone program

$$\begin{array}{ll} \text{minimize} & t \\ \text{subject to} & (-t, q, p) \in \mathcal{E} \quad (\Leftrightarrow \exp(-t/p) \leq q/p) \end{array}$$

with variable t

SDP representations

Nesterov, Nemirovsky, and others have worked out SDP representations for many functions, *e.g.*,

- x^p , $p \geq 1$ rational
- $-(\det X)^{1/n}$
- $\sum_{i=1}^k \lambda_i(X)$ ($X = X^T$)
- $\|X\| = \sigma_1(X)$ ($X \in \mathbf{R}^{m \times n}$)
- $\|X\|_* = \sum_i \sigma_i(X)$ ($X \in \mathbf{R}^{m \times n}$)

some of these representations are not obvious . . .

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Transforming to cone program: Example

- example: ℓ_1 -regularized least-norm problem

$$\text{minimize } \|Ax - b\|_2 + \lambda\|x\|_1$$

with variable $x \in \mathbf{R}^n$ (convex, but not a cone program . . .)

- introduce new variables $t \in \mathbf{R}$, $z \in \mathbf{R}^m$, x_+ , $x_- \in \mathbf{R}^n$:

$$\begin{aligned} \text{minimize } & t + \lambda(\mathbf{1}^T x_+ + \mathbf{1}^T x_-) \\ \text{subject to } & \|z\|_2 \leq t, \quad x_+ \succeq 0, \quad x_- \succeq 0 \\ & z = Ax - b, \quad x = x_+ - x_- \end{aligned}$$

. . . a cone program with $x \in \mathbf{R}^n$, $(x_+, x_-) \in \mathbf{R}_+^{2n}$, $(z, t) \in \mathcal{L}^m$

- optimal x for cone program is optimal for original problem

Transforming to cone program: General case

- start with convex optimization problem \mathcal{P}_0
- find sequence of transformations that yields cone program \mathcal{P}_K

$$\mathcal{P}_0 \rightarrow \mathcal{P}_1 \rightarrow \cdots \rightarrow \mathcal{P}_K$$

- solve \mathcal{P}_K efficiently
- transform solution of \mathcal{P}_K back to solution of original problem \mathcal{P}_0

Problem transformations

(a.k.a. re-writing methods, reductions)

- there are many such problem transformations, some obvious, many not
- idea goes back at least to 1940s (Dantzig, for LP)
- standard optimization curriculum covers transformation ‘tricks’
(to be carried out by hand, *i.e.*, graduate student)

Convex calculus rules and problem transformations

- for most of the convex calculus rules, there is an associated problem transformation that ‘undoes’ the rule
- example: suppose $\max\{f_1(x), f_2(x)\}$ appears as term in objective or constraint function, with f_1, f_2 convex
- problem transformation:
 - replace $\max\{f_1(x), f_2(x)\}$ with a new variable t
 - add new (convex) constraints $f_1(x) \leq t, f_2(x) \leq t$
- yields equivalent problem
(solution x of new problem is solution of original problem)

General composition rule and problem transformations

- suppose we encounter expression $h(f_1(x), \dots, f_k(x))$, convex by general composition rule
- problem transformation:
 - replace expression with new variable t
 - add new variables s_1, \dots, s_k and (convex) constraints

$$\begin{array}{ll} f_i(x) \leq s_i & f_i \text{ convex} \\ f_i(x) \geq s_i & f_i \text{ concave} \\ f_i(x) = s_i & f_i \text{ affine} \end{array}$$

- add (convex) constraint $h(s_1, \dots, s_k) \leq t$
- yields equivalent problem

Cone representations and problem transformations

- suppose f has cone representation

$$\begin{array}{ll} \text{minimize} & c^T x + d^T y + e \\ \text{subject to} & A \begin{bmatrix} x \\ y \end{bmatrix} = b, \quad \begin{bmatrix} x \\ y \end{bmatrix} \in \mathcal{K} \end{array}$$

- when we encounter $f(x)$, we
 - add new variable y , and constraints above
 - replace $f(x)$ with $c^T x + d^T y + e$
- yields equivalent problem

Transforming from DCP to cone problem

- start with convex program in DCP form with cone-representable functions
- parse tree gives
 - proof of convexity
 - set of transformations that reduce original problem to an equivalent cone program

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Parser/solver

- specify convex problem in natural (DCP) form
 - declare optimization variables
 - form convex objective and constraints using functions from a library, general composition rule
 - library functions defined by cone representations
- parser/solver
 - parses problem description (certifying convexity)
 - automatically transforms to equivalent cone problem
 - solves cone problem
 - transforms back to get solution of original problem

Parser/solver

- implemented using object-oriented methods and/or compiler-compilers
- huge gain in productivity
 - rapid prototyping
 - teaching
 - trying out research ideas
- transformation can introduce many new variables and constraints, but performance penalty small if sparsity is exploited

History

- general purpose optimization modeling systems AMPL, GAMS (1970s)
- systems for SDPs/LMIs (1990s): `sdpsol` (Wu, Boyd), `lmilab` (Gahinet, Nemirovsky), `lmitool` (El Ghaoui)
- `yalmip` (Löfberg 2000–)
- automated convexity checking (Crusius PhD thesis 2002)
- disciplined convex programming (Grant, Boyd, Ye 2004)
- `cvx` (Grant, Boyd 2005)
- `cvxopt` (Dahl, Vandenberghe 2005)
- `ggplab` (Mutapcic, Koh, et al 2006)

Example: CVX

- parser/solver written in Matlab (M. Grant)
- convex problem, with variable $x \in \mathbf{R}^n$; A, b, λ, F, g constants

$$\begin{aligned} & \text{minimize} && \|Ax - b\|_2 + \lambda\|x\|_1 \\ & \text{subject to} && Fx \leq g \end{aligned}$$

- CVX specification:

```
cvx_begin
    variable x(n)          % declare vector variable
    minimize (norm(A*x-b,2) + lambda*norm(x,1))
    subject to F*x <= g
cvx_end
```

when CVX processes this specification, it

- parses it (which verifies convexity of problem)
- generates equivalent cone problem (here, an SOCP)
- solves it using SDPT3 or SeDuMi
- transforms solution back to original problem

the CVX code is easy to read, understand, modify

The same example, transformed by 'hand'

transform problem to SOCP, call SeDuMi, reconstruct solution:

```
% Set up big matrices.
[m,n] = size(A); [p,n] = size(F);
AA = [speye(n), -speye(n), speye(n), sparse(n,p+m+1); ...
      F, sparse(p,2*n), speye(p), sparse(p,m+1); ...
      A, sparse(m,2*n+p), speye(m), sparse(m,1)];
bb = [zeros(n,1); g; b];
cc = [zeros(n,1); gamma*ones(2*n,1); zeros(m+p,1); 1];
K.f = m; K.l = 2*n+p; K.q = m + 1;      % specify cone
xx = sedumi(AA, bb, cc, K);              % solve SOCP
x = x(1:n);                             % extract solution
```

Some functions in the CVX library

function	meaning	attributes
<code>norm(x, p)</code>	$\ x\ _p, p \geq 1$	cvx
<code>square(x)</code>	x^2	cvx
<code>square_pos(x)</code>	$(x_+)^2$	cvx, nondecr
<code>pos(x)</code>	x_+	cvx, nondecr
<code>sum_largest(x, k)</code>	$x_{[1]} + \dots + x_{[k]}$	cvx, nondecr
<code>sqrt(x)</code>	$\sqrt{x}, x \geq 0$	ccv, nondecr
<code>inv_pos(x)</code>	$1/x, x > 0$	cvx, nonincr
<code>max(x)</code>	$\max\{x_1, \dots, x_n\}$	cvx, nondecr
<code>quad_over_lin(x, y)</code>	$x^2/y, y > 0$	cvx, nonincr in y
<code>lambda_max(X)</code>	$\lambda_{\max}(X), X = X^T$	cvx
<code>huber(x)</code>	$\begin{cases} x^2, & x \leq 1 \\ 2 x - 1, & x > 1 \end{cases}$	cvx

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Solving specific problems

in developing a custom solver for a specific application, we can

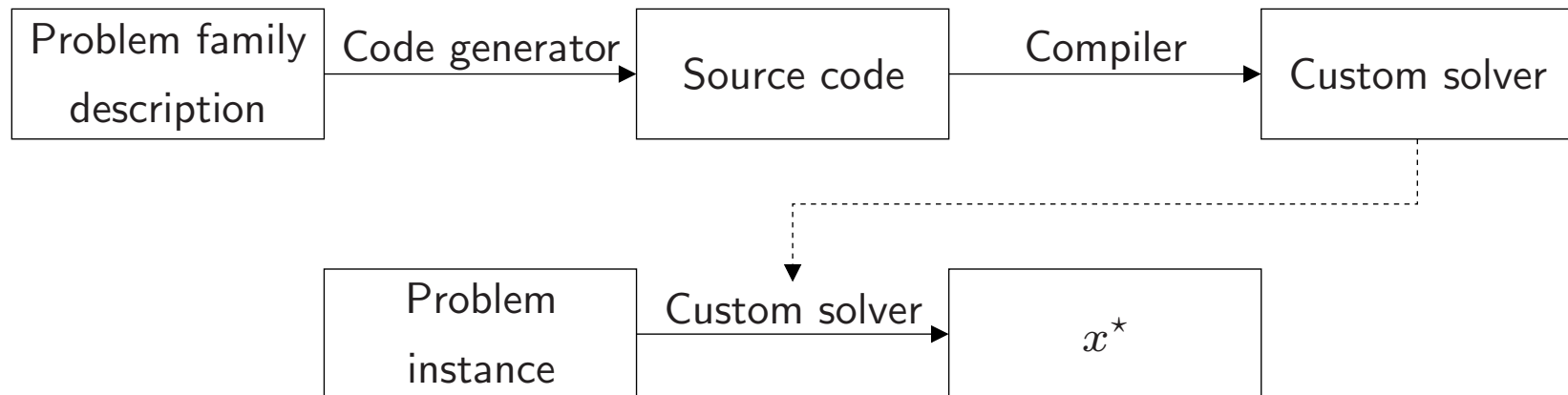
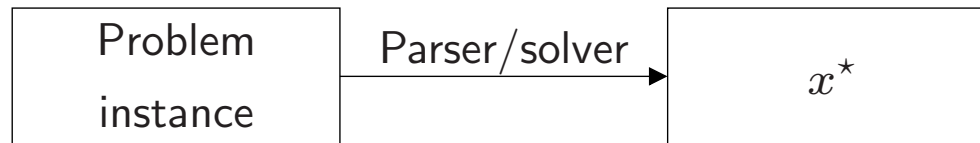
- optimize transformation to cone form
- exploit structure very efficiently
- determine ordering, memory allocation beforehand
- cut corners in algorithm, *e.g.*, terminate early
- use warm start

to get **very fast** solver

Convex optimization solver generation

- specify convex problem **family** via (an extension of) DCP
 - declare optimization variables, *parameters*
 - form convex objective and constraints from library of functions, using general composition rule
 - *parameter constraints* used in parsing
- code generator
 - analyzes problem structure (dimensions, sparsity, . . .)
 - chooses elimination ordering
 - generates solver code for specific problem family
- idea:
 - spend (perhaps much) time generating code
 - save (hopefully much) time solving problem instances

Parser/solver vs. code generation



Example: CVXMOD

- written in Python (J. Mattingley)
- QP family, with variable $x \in \mathbf{R}^n$, parameters P, q, g, h

$$\begin{aligned} & \text{minimize} && x^T P x + q^T x \\ & \text{subject to} && G x \leq h, \quad A x = b \end{aligned}$$

- CVXMOD specification:

```
A = matrix(...); b = matrix(...)
P = param('P', n, n, psd=True); q = param('q', n)
G = param('G', m, n); h = param('h', m)
x = optvar('x', n)
qpfam = problem(minimize(quadform(x, P) + tp(q)*x),
                [G*x <= h, A*x == b])
```

cvxmod code generation

- generate solver for problem family `qpfam` with

```
qpfam.codegen()
```

- output includes `qpfam/solver.c` and ancillary files
- solve instance with (C function call)

```
status = solver(params, vars, work);
```

cvxmod code generator

(preliminary implementation)

- handles problems transformable to QP
- primal-dual interior-point method
- direct LDL^T factorization of KKT matrix
- (slow) method to determine variable ordering (at code generation time)
- explicit factorization code generated

Sample solve times for CVXMOD generated code

problem family	vars	constrs	SDPT3 (ms)	cvxmod (ms)
control1	140	190	250	0.4
control2	360	1080	1400	2.0
control3	1110	3180	3400	13.2
order_exec	20	41	490	0.05
net_utility	50	150	130	0.23
actuator	50	106	300	0.17
robust_kalman	95	45	120	0.12

Conclusions

- DCP is a formalization of constructive convex analysis
 - simple method to certify problem as convex
 - useful as language for describing convex problems
- parser/solvers make rapid prototyping easy
- new code generation methods yield solvers that
 - are extremely fast
 - can be embedded in real-time applications

References

- *Disciplined Convex Programming* (Grant, Boyd, Ye)
- *Graph Implementations for Nonsmooth Convex Programs* (Grant, Boyd)
- *Automatic Code Generation for Real-Time Convex Optimization* (Mattingley, Boyd)
- CVX (Grant, Boyd)
- CVXMOD (Mattingley, Boyd)

all available on-line, but CVXMOD code gen not yet ready for prime-time