

Weighted Sum Rate Optimization for Cognitive Radio MIMO Broadcast Channels

Lan Zhang[†], Yan Xin^{†*}, and Ying-Chang Liang[‡]

Abstract

In this paper, we consider a cognitive radio (CR) network in which the unlicensed (secondary) users are allowed to concurrently access the spectrum allocated to the licensed (primary) users provided that their interference to the primary users (PUs) satisfies certain constraints. We study a weighted sum rate maximization problem for the secondary user (SU) multiple input multiple output (MIMO) broadcast channel (BC), in which the SUs have not only the sum power constraint but also interference constraints. We transform this multi-constraint maximization problem into its equivalent form, which involves a single constraint with multiple auxiliary variables. Fixing these multiple auxiliary variables, we propose a duality result for the equivalent problem. Our duality result can solve the optimization problem for MIMO-BC with multiple linear constraints, and thus can be viewed as a generalization of the conventional results, which rely crucially on a single sum power constraint. Furthermore, we develop an efficient sub-gradient based iterative algorithm to solve the equivalent problem and show that the developed algorithm converges to a globally optimal solution. Simulation results are further provided to corroborate the effectiveness of the proposed algorithm.

Index Terms

Beamforming, broadcast channel, cognitive radio, MIMO, power allocation, sum rate maximization.

[†]The authors are with the Department of Electrical and Computer Engineering, National University of Singapore, Singapore 118622 (email: zhanglan@nus.edu.sg; elexy@nus.edu.sg).

[‡]The author is with Institute of Infocomm Research, 21 Heng Mui Keng Terrace, Singapore 119613 (email: ycliang@i2r.a-star.edu.sg). The corresponding author of the paper is Dr. Ying-Chang Liang (Tel. no: (65) 6874-8225).

I. INTRODUCTION

Cognitive radio (CR), as a promising technology to advocate efficient use of radio spectrum, has been a topic of increasing research interest in recent years [1]–[6]. CR allows an unlicensed (secondary) user to *opportunistically* or *concurrently* access the spectrum initially allocated to the licensed (primary) users provided that certain prescribed constraints are satisfied, thus having a potential to improve spectral utilization efficiency. In this paper, we study a weighted sum rate maximization problem for the secondary user (SU) multiple input multiple output (MIMO) broadcast channel (BC) in a concurrent CR network, in which the SUs have not only the sum power constraint but also interference constraints.

A. System Model and Problem Formulation

With reference to Fig. 1, we consider the K -SU MIMO-BC with N_t transmit antennas and N_r receive antennas in a CR network, where the K SUs share the same spectrum with a single primary user (PU) equipped with one transmitter and one receiver¹. The transmit-receive signal model from the BS to the i th SU denoted by SU_i , for $i = 1, \dots, K$, can be expressed as

$$\mathbf{y}_i = \mathbf{H}_i \mathbf{x} + \mathbf{z}_i, \quad (1)$$

where \mathbf{y}_i is the $N_r \times 1$ received signal vector, \mathbf{H}_i is the $N_r \times N_t$ channel matrix from the BS to the SU_i , \mathbf{x} is the $N_t \times 1$ transmitted signal vector, and \mathbf{z}_i is the $N_r \times 1$ Gaussian noise vector with entries being independent identically distributed random variables (RVs) with means zero and variances σ^2 . Consider \mathbf{h}_o as the $N_t \times 1$ channel gain vector between the transmitters of the BS and the PU. We further assume that \mathbf{H}_i for $i = 1, \dots, K$, and \mathbf{h}_o remain constant during a transmission block and change independently from block to block, and \mathbf{H}_i for $i = 1, \dots, K$, and \mathbf{h}_o are perfectly known to the BS of the SUs and SU_i . To acquire channel matrices \mathbf{H}_i and channel vector \mathbf{h}_o at the BS, transmission protocols need to incorporate certain cooperation among the PU, the SUs, and the BS of the SUs. As an example, the BS need transmit pilot symbols to allow the SUs and PU to obtain respective estimates of channel matrices \mathbf{H}_i and channel vector \mathbf{h}_o , and such estimates are needed to be reliably transmitted back to the BS via feedback channels.

¹Except for explicitly stated, we restrict our attention to a single PU case in the rest of this paper for convenience of description. The results derived for the single PU case can be readily extended to the multiple PU case, which will be discussed in Remark 3.

We next consider the weighted sum rate maximization problem for the K -SU MIMO-BC in a CR network, simply called the CR MIMO-BC sum rate maximization problem, which mathematically can be formulated as

Problem 1 (Main Problem):

$$\begin{aligned} \max_{\{\mathbf{Q}_i^b\}_{i=1}^K: \mathbf{Q}_i^b \succeq 0} & \sum_{i=1}^K w_i r_i^b \\ \text{subject to} & \sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o \leq P_t, \text{ and } \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) \leq P_u, \end{aligned} \quad (2)$$

where r_i^b is the rate achieved by SU $_i$, w_i is the weight of SU $_i$, \mathbf{Q}_i^b denotes the $N_t \times N_t$ transmit signal covariance matrix for SU $_i$, $\mathbf{Q}_i^b \succeq 0$ indicates that \mathbf{Q}_i^b is a semidefinite matrix, P_t denotes the interference threshold of the PU, and P_u denotes the sum power constraint at the BS. Compared with the same under a non-CR setting, the key difference is that in addition to the sum power constraint, an interference constraint is applied to the SUs in the CR MIMO-BC, i.e., the total received interference power $\sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o$ at the PU is below the threshold P_t .

B. Related Work

The present paper is primarily motivated by the previous work on the information-theoretic study of the MIMO-BC under a non-CR setting. It has been shown in [7] [8] [9] that under a single sum power constraint, the sum rate capacity (equally weighted) of the non-CR MIMO-BC is achieved by the dirty paper coding (DPC) scheme. Furthermore, the paper [10] shows that the rate region achieved by the DPC scheme is indeed the capacity region of such a channel. However, the power allocation and beamforming strategy that maximizes the sum rate capacity have not been investigated in these papers. Under a single sum power constraint, several efficient algorithms have been developed to obtain the optimal solutions for the MIMO-BC sum rate problems, including both the equally weighted case [11], [12] and the weighted case [13], [14]. The development of these algorithms is crucially relied on the BC-MAC duality result [15], which is only valid for the problem with a single sum power constraint. Due to this limitation, the aforementioned algorithms are not applicable to the problem with multiple linear constraints, which is the case of interest in this paper. To cope with a problem with multiple linear constraints, the paper [16] proposed a novel minimax BC-MAC duality. Based on this duality, an interior point algorithm was proposed to solve the MIMO-BC weighted sum rate problem with per-antenna power constraints [17].

C. Contribution

Throughout this paper, we consider the CR MIMO-BC weighted sum rate maximization problem as defined in Problem 1, where the transmit power is subject to general multiple linear constraints instead of the sum power constraint or per-antenna power constraints. As the main contribution of this paper, our solution is summarized in the following.

- 1) We prove that in the CR MIMO-BC, the multi-constraint weighted sum rate maximization problem (Problem 1) is equivalent to a single-constraint weighted sum rate maximization problem with multiple auxiliary variables.
- 2) For the equivalent problem, we establish a duality between the MIMO-BC and a dual MIMO-MAC when the multiple auxiliary variables are fixed as constant. The new duality result is applicable to the MIMO-BC weighted sum rate problem with general multiple linear constraints, and therefore can be viewed as a generalization of the conventional MIMO BC-MAC duality [15] [7] [8] that restricts to a single sum power constraint. Compared with the minimax BC-MAC duality [16], our duality result is given in a simpler form.
- 3) We develop an efficient algorithm, which employs a primal dual approach [18], to solve the dual MIMO-MAC weighted sum rate maximization problem. It should be noted that the dual MIMO-MAC weighted sum rate problem cannot be solved by using existing iterative water-filling based algorithms [11], [12]. Furthermore, we propose a sub-gradient based iterative algorithm to solve the main problem of the paper, Problem 1, and show that the proposed algorithm converges to the globally optimal solution. Unlike the interior point based algorithm proposed in [17], our proposed algorithm relies on the Karush-Kuhn-Tucker (KKT) conditions.

D. Organization and Notation

The rest of this paper is organized as follows. In Section II, we transform the CR MIMO-BC weighted sum rate maximization problem (Problem 1) into its equivalent form, and introduce a MAC-BC duality between a MIMO-BC and a dual MIMO-MAC. Section III presents an primal dual method based iterative to solve the dual MIMO-MAC weighted sum rate problem. Section IV develops the complete algorithm to solve the CR MIMO-BC weighted sum rate maximization problem. Section V provides several simulation examples. Finally, Section VI concludes the paper.

The following notations are used in this paper. Bold upper and lower case letter denote matrices and vectors, respectively; $(\cdot)^\dagger$ and $(\cdot)^T$ denote the conjugate transpose and transpose, respectively; \mathbf{I}_M denotes

an $M \times M$ identity matrix; $\text{tr}(\cdot)$ denotes the trace of a matrix, and $[x]^+$ denotes $\max(x, 0)$; $(\cdot)^b$ and $(\cdot)^m$ denote the quantities associated with a broadcast channel and a multiple access channel, respectively; $E[\cdot]$ denotes the expectation operator.

II. EQUIVALENCE AND DUALITY

Evidently, the MIMO-BC weighted sum rate maximization problem under either a non-CR or a CR setting is a non-convex optimization problem and is difficult to solve directly. Under a single sum power constraint, the weighted sum rate problem for MIMO BC can be transformed to its dual MIMO MAC problem, which is convex and can be solved in an efficient manner [11]. In the CR setting, the problem (Problem 1) has not only a sum power constraint but also an interference constraint. The imposed multiple constraints render difficulty to formulate an efficiently solvable dual problem. To overcome the difficulty, we first transform this multi-constrained weighted sum rate problem (Problem 1) into its equivalent problem which has a single constraint with multiple auxiliary variables, and next develop a duality between a MIMO-BC and a dual MIMO-MAC in the case where the multiple auxiliary variables are fixed.

A. An Equivalent MIMO-BC Weighted Sum Rate Problem

In the following, we present an equivalent form of Problem 1 (see Appendix A for the proof).

Proposition 1: Problem 1 shares the same optimal solution with

Problem 2 (Equivalent Problem):

$$\min_{q_t \geq 0, q_u \geq 0} \max_{\{\mathbf{Q}_i^b\}_{i=1}^K, \mathbf{Q}_i^b \geq 0} \sum_{i=1}^K w_i r_i^b \quad (3)$$

$$\text{subject to } q_t \left(\sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o - P_t \right) + q_u \left(\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P_u \right) \leq 0, \quad (4)$$

where q_t and q_u are the auxiliary dual variables for the respective interference constraint and sum power constraint.

Finding an efficiently solvable dual problem for Problem 2 directly is still difficult. However, as we show later, when q_t and q_u are fixed as constants, Problem 2 reduces to a simplified form, which we can solve by applying the following duality result.

B. CR MIMO BC-MAC Duality

For fixed q_t and q_u , Problem 2 reduces to the following form

Problem 3 (CR MIMO-BC):

$$\max_{\{\mathbf{Q}_i^b\}_{i=1}^K: \mathbf{Q}_i^b \succeq 0} \sum_{i=1}^K w_i r_i^b \quad (5)$$

$$\text{subject to } q_t \sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o + q_u \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) \leq P, \quad (6)$$

where $P := q_t P_t + q_u P_u$. Since q_t and q_u are fixed, P is a constant in Problem 3. The constraint (6) is not a single sum power constraint, and thus the duality result established in [11] is not applicable to Problem 3. Therefore, we formulate the following new dual MAC problem.

Proposition 2: The dual MAC problem of Problem 3 is

Problem 4 (CR MIMO-MAC):

$$\max_{\{\mathbf{Q}_i^m\}_{i=1}^K: \mathbf{Q}_i^m \succeq 0} \sum_{i=1}^K w_i r_i^m \quad (7)$$

$$\text{subject to } \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) \sigma^2 \leq P, \quad (8)$$

where r_i^m is the rate achieved by the i th user of the dual MAC, and \mathbf{Q}_i^m is the transmit signal covariance matrix of the i th user, and the noise covariance at the BS is $q_t \mathbf{h}_o \mathbf{h}_o^H + q_u \mathbf{I}_{N_t}$.

Remark 1: According to Proposition 2, for fixed q_t and q_u , the optimal weighted sum rate of the dual MAC is equal to the optimal weighted sum rate of the primal BC. From the formulation perspective, this duality result is quite similar to the conventional duality in [15] [7] [8]. However, as shown in Fig. 2, one thing needs to highlight is that the noise covariance matrix of the dual MAC is a function of the auxiliary variable q_t and q_u , instead of the identity matrix [8]. This difference comes from the constraint (6), which is not a sum power constraint as in [8]. Note that when $q_t = 0$, the duality result reduces to the conventional BC-MAC duality in [8]. Compared with the minimax duality in [17], our duality result has a simpler format.

As illustrated in Fig. 2, Proposition 2 describes a weighted sum rate maximization problem for a dual MIMO-MAC. To prove the proposition, we first examine the relation between the *signal to interference plus noise ratio* (SINR) regions of the MIMO-BC and the dual MIMO-MAC. Based on this relation, we will show that the achievable rate regions of the MIMO-BC and the dual MIMO-MAC are the same.

In the sequel, we first describe the definition of the SINR for the MIMO-BC. It has been shown in [10] that the DPC is a capacity achieving scheme. Each set of the transmit covariance matrix determined by DPC scheme defines a set of transmit and receive beamforming vectors, and each pair of these transmit

and receive beamforming vectors forms a data stream. In a beamforming perspective, the BS transmitter have $N_t \times K$ beamformers, $\mathbf{u}_{i,j}$, for $i = 1, \dots, K$, and $j = 1, \dots, N_t$. Therefore, the transmit signal can be represented as

$$\mathbf{x} = \sum_{i=1}^K \sum_{j=1}^{N_t} x_{i,j} \mathbf{u}_{i,j},$$

where $x_{i,j}$ is a scalar representing the data stream transmitted in this beamformer, and $E[x_{i,j}^2] = p_{i,j}$ denotes the power allocated to this beamformer. At SU_i , the receive beamformer corresponding to $\mathbf{u}_{i,j}$ is denoted by $\mathbf{v}_{i,j}$. The transmit beamformer $\mathbf{u}_{i,j}$ and the power $p_{i,j}$ can be obtained via the eigenvalue decomposition of \mathbf{Q}_i^b , i.e., $\mathbf{Q}_i^b = \mathbf{U}_i^\dagger \mathbf{P}_i \mathbf{U}_i$, where \mathbf{U}_i is a unitary matrix, and \mathbf{P}_i is a diagonal matrix. The transmit beamformer $\mathbf{u}_{i,j}$ is the j th column of \mathbf{U}_i , and $p_{i,j}$ is the j th diagonal entry of \mathbf{P}_i . With these notations, we express the $\text{SINR}_{i,j}^b$ as

$$\text{SINR}_{i,j}^b = \frac{p_{i,j} |\mathbf{u}_{i,j}^\dagger \mathbf{H}_i^\dagger \mathbf{v}_{i,j}|^2}{\sum_{k=i+1}^K \sum_{l=1}^{N_r} p_{k,l} |\mathbf{u}_{k,l}^\dagger \mathbf{H}_i^\dagger \mathbf{v}_{i,j}|^2 + \sum_{l=j+1}^{N_r} p_{i,l} |\mathbf{u}_{i,l}^\dagger \mathbf{H}_i^\dagger \mathbf{v}_{i,j}|^2 + \sigma^2}. \quad (9)$$

It can be observed from (9) that the DPC scheme is applied. This can be interpreted as follows. The signal from SU_1 is first encoded with the signals from other SUs being treated as interference. The signal from SU_2 is next encoded by using the DPC scheme. Signals from the other SUs will be encoded sequentially in a similar manner. For the data streams within SU_i , the data stream 1 is also encoded first while the other data streams are treated as the interference. The data stream 2 is encoded next. In a similar manner, the other data streams will be sequentially encoded. The encoding order is assumed to be arbitrary at this moment, and the optimal encoding order of Problem 2 will be discussed in Section III.

To explore the relation of the SINR regions of the dual MAC and the BC, we formulate a following optimization problem

$$\min_{\{\mathbf{Q}_i^b\}_{i=1}^K: \mathbf{Q}_i^b \succeq 0} q_t \sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o + q_u \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P, \text{ subject to: } \text{SINR}_{i,j}^b \geq \gamma_{i,j}, \quad (10)$$

where $\gamma_{i,j}$ denotes the SINR threshold of the j th data stream within the SU_i for the BC. Note that the objective function in (10) is a function of signal covariance matrices and the constraints are SINR constraints for the K -SU MIMO-BC.

It has been shown in [17] and [19] that the *non-convex* BC sum power minimization problem under the SINR constraints can be solved efficiently via its dual MAC problem, which is a convex optimization problem. By following a similar line of thinking, the problem in (10) can be efficiently solved via its dual MAC problem. Similar to the primal MIMO-BC, the dual MIMO-MAC depicted in Fig. 2 consists

of K users each with N_r transmit antennas, and one BS with N_t receive antennas. By transposing the channel matrix and interchanging the input and output signals, we obtain the dual MIMO-MAC from the primal MIMO-BC. For the covariance matrices \mathbf{Q}_i^m of the dual MIMO-MAC, we apply the eigenvalue decomposition,

$$\mathbf{Q}_i^m = \mathbf{V}_i \mathbf{\Lambda}_i \mathbf{V}_i^\dagger = \sum_{j=1}^{N_r} q_{i,j} \mathbf{v}_{i,j} \mathbf{v}_{i,j}^\dagger, \quad (11)$$

where $\mathbf{v}_{i,j}$ is the j th column of \mathbf{V}_i , and $q_{i,j}$ is the j th diagonal entry of $\mathbf{\Lambda}_i$. For user i , $\mathbf{v}_{i,j}$ is the transmit beamforming vector of the j th data stream, the power allocated to the j th data stream equals $q_{i,j}$, and the receive beamforming vector of the j th data stream at the BS is $\mathbf{u}_{i,j}$. The SINR of the dual MIMO-MAC is given by

$$\text{SINR}_{i,j}^m = \frac{q_{i,j} |\mathbf{u}_{i,j}^\dagger \mathbf{H}_i^\dagger \mathbf{v}_{i,j}|^2}{\mathbf{u}_{i,j}^\dagger \left(\sum_{k=1}^{i-1} \sum_{l=1}^{N_r} q_{k,l} \mathbf{H}_k^\dagger \mathbf{v}_{k,l} \mathbf{v}_{k,l}^\dagger \mathbf{H}_k + \sum_{l=1}^{j-1} q_{i,l} \mathbf{H}_i^\dagger \mathbf{v}_{i,l} \mathbf{v}_{i,l}^\dagger \mathbf{H}_i + \mathbf{R}_w \right) \mathbf{u}_{i,j}}, \quad (12)$$

where $\mathbf{R}_w := q_t \mathbf{R}_o + q_u \mathbf{I}_{N_t}$ is the noise covariance matrix of the MIMO-MAC with $\mathbf{R}_o := \mathbf{h}_o \mathbf{h}_o^\dagger$. In the dual MIMO-MAC, \mathbf{R}_w depends on q_t and q_u defined in (10) whereas the noise covariance matrix in the primal MIMO-BC is an identity matrix. It can be observed from (12) that the successive interference cancelation (SIC) scheme is used in this dual MIMO-MAC, and the decoding order is the reverse encoding order of the primal BC. The signal from SU_K is first decoded with the signals from other users being treated as interference. After decoded at the BS, the signals from SU_K will be subtracted from the received signal. The signal from SU_{K-1} is next decoded, and so on. Again, the data streams within a SU can be decoded in a sequential manner.

For the dual MIMO-MAC, we consider the following minimization problem similar to the problem (10)

$$\min_{\{\mathbf{Q}_i^m\}_{i=1}^K: \mathbf{Q}_i^m \succeq 0} \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) \sigma^2 - P, \quad \text{subject to: } \text{SINR}_{i,j}^m \geq \gamma_{i,j}. \quad (13)$$

The following proposition describes the relation between the problems (10) and (13).

Proposition 3: For fixed q_t and q_u , the MIMO-MAC problem (13) is dual to the MIMO-BC problem (10).

Proof: The constraints in (10) can be rewritten as

$$\frac{p_{i,j} |\mathbf{u}_{i,j}^\dagger \mathbf{H}_i^\dagger \mathbf{v}_{i,j}|^2}{\gamma_{i,j}} \geq \sum_{k=i+1}^K \sum_{l=1}^{N_r} p_{k,l} |\mathbf{u}_{k,l}^\dagger \mathbf{H}_k^\dagger \mathbf{v}_{k,l}|^2 + \sum_{l=j+1}^{N_r} p_{i,l} |\mathbf{u}_{i,l}^\dagger \mathbf{H}_i^\dagger \mathbf{v}_{i,l}|^2 + \sigma^2. \quad (14)$$

Thus, the Lagrangian function of the problem (10) is

$$\begin{aligned}
& L_1(\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b, \lambda_{i,j}) \\
&= q_t \sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o + q_u \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P - \sum_{i=1}^K \sum_{j=1}^{N_r} \lambda_{i,j} \left(\frac{p_{i,j} |\mathbf{u}_{i,j}^\dagger \mathbf{H}_i^\dagger \mathbf{v}_{i,j}|^2}{\gamma_{i,j}} \right. \\
&\quad \left. - \sum_{k=i+1}^K \sum_{l=1}^{N_r} p_{k,l} |\mathbf{u}_{k,l}^\dagger \mathbf{H}_i^\dagger \mathbf{v}_{i,j}|^2 - \sum_{l=j+1}^{N_r} p_{i,l} |\mathbf{u}_{i,l}^\dagger \mathbf{H}_i^\dagger \mathbf{v}_{i,j}|^2 - \sigma^2 \right) \tag{15}
\end{aligned}$$

$$\begin{aligned}
&= \sum_{i=1}^K \sum_{j=1}^{N_r} \lambda_{i,j} \sigma^2 - P - \sum_{i=1}^K \sum_{j=1}^{N_r} p_{i,j} \mathbf{u}_{i,j}^\dagger \left(\frac{\lambda_{i,j} \mathbf{H}_i^\dagger \mathbf{v}_{i,j} \mathbf{v}_{i,j}^\dagger \mathbf{H}_i}{\gamma_{i,j}} \right. \\
&\quad \left. - \sum_{k=1}^{i-1} \sum_{l=1}^{N_r} \lambda_{k,l} \mathbf{H}_k^\dagger \mathbf{v}_{k,l} \mathbf{v}_{k,l}^\dagger \mathbf{H}_k - \sum_{l=1}^{j-1} \lambda_{i,l} \mathbf{H}_i^\dagger \mathbf{v}_{i,l} \mathbf{v}_{i,l}^\dagger \mathbf{H}_i - \mathbf{R}_w \right) \mathbf{u}_{i,j}, \tag{16}
\end{aligned}$$

where $\lambda_{i,j}$ is the Lagrangian multiplier. Eq. (16) is obtained by applying the eigenvalue decomposition to \mathbf{Q}_i^b and rearranging the terms in (15). The optimal objective value of (10) is

$$\max_{\lambda_{i,j}} \min_{\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b} L_1(\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b, \lambda_{i,j}). \tag{17}$$

On the other hand, the Lagrangian function of the problem (13) is

$$\begin{aligned}
L_2(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m, \delta_{i,j}) &= \sum_{i=1}^K \sum_{j=1}^{N_r} q_{i,j} \sigma^2 - P - \sum_{i=1}^K \sum_{j=1}^{N_r} \delta_{i,j} \mathbf{u}_{i,j}^\dagger \left(\frac{q_{i,j} \mathbf{H}_i^\dagger \mathbf{v}_{i,j} \mathbf{v}_{i,j}^\dagger \mathbf{H}_i}{\gamma_{i,j}} \right. \\
&\quad \left. - \sum_{k=1}^{i-1} \sum_{l=1}^{N_r} q_{k,l} \mathbf{H}_k^\dagger \mathbf{v}_{k,l} \mathbf{v}_{k,l}^\dagger \mathbf{H}_k - \sum_{l=1}^{j-1} q_{i,l} \mathbf{H}_i^\dagger \mathbf{v}_{i,l} \mathbf{v}_{i,l}^\dagger \mathbf{H}_i - \mathbf{R}_w \right) \mathbf{u}_{i,j}, \tag{18}
\end{aligned}$$

where $\delta_{i,j}$ is the Lagrangian multiplier. Eq. (18) is also obtained by applying eigenvalue decomposition to \mathbf{Q}_i^m . The optimal objective value of (13) is

$$\max_{\delta_{i,j}} \min_{\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m} L_2(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m, \delta_{i,j}). \tag{19}$$

Note that if we choose $q_{i,j} = \lambda_{i,j}$, $\delta_{i,j} = p_{i,j}$, and the same beamforming vectors $\mathbf{u}_{i,j}$ and $\mathbf{v}_{i,j}$ for both problems, (16) and (18) become identical. This means that the optimal solutions of (17) and (19) are the same. \blacksquare

Proposition 3 implies that under the SINR constraints, the problems (10) and (13) can achieve the same objective value, which is a function of the transmit signal covariance matrices. On the other hand, under the corresponding constraints on the signal covariance matrix, the achievable SINR regions of the

MIMO-BC and its dual MIMO-MAC are the same. Mathematically, we define the respective achievable SINR regions for the primal MIMO-BC and the dual MIMO-MAC as follows.

Definition 1: A SINR vector $\gamma = (\gamma_{1,1}, \dots, \gamma_{1,N_t}, \dots, \gamma_{K,N_t})$ is said to be achievable for the primal BC if and only if there exists a set of $\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b$ such that $q_t \sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o + q_u \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P \leq C$ for a constant C and the corresponding $\text{SINR}_{i,j}^b \geq \gamma_{i,j}$. An achievable BC SINR region denoted by \mathcal{R}_{BC} , is a set containing all the BC achievable γ .

Definition 2: A SINR vector $\gamma = (\gamma_{1,1}, \dots, \gamma_{1,N_t}, \dots, \gamma_{K,N_t})$ is said to be achievable for the dual MAC if and only if there exists a set of $\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m$ such that $\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) \sigma^2 - P \leq C$ for a constant C and the corresponding $\text{SINR}_{i,j}^m \geq \gamma_{i,j}$. An achievable MAC SINR region denoted by \mathcal{R}_{MAC} , is a set containing all the MAC achievable γ .

In the following corollary, we will show $\mathcal{R}_{\text{MAC}} = \mathcal{R}_{\text{BC}}$.

Corollary 1: For fixed q_t and q_u , and a constant C , the MIMO-BC under the constraint $q_t \sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o + q_u \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P \leq C$ and the dual MIMO-MAC under the constraint $\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) \sigma^2 - P \leq C$ achieve the same SINR region.

Proof: For any $\gamma \in \mathcal{R}_{\text{MAC}}$, by Definition 2, there exists a set of $\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m$ such that $\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) \sigma^2 - P \leq C$ and the corresponding $\text{SINR}_{i,j}^m \geq \gamma_{i,j}$. It can be readily concluded from Proposition 3 that there exists a set of $\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b$ such that $q_t \sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o + q_u \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P \leq C$ and the corresponding $\text{SINR}_{i,j}^b \geq \gamma_{i,j}$. This implies $\gamma \in \mathcal{R}_{\text{BC}}$. Since γ is an arbitrary element in \mathcal{R}_{MAC} , we have $\mathcal{R}_{\text{MAC}} \subseteq \mathcal{R}_{\text{BC}}$. In a similar manner, we have $\mathcal{R}_{\text{BC}} \subseteq \mathcal{R}_{\text{MAC}}$. The proof follows. ■

We are now in the position to prove Proposition 2.

Proof of Proposition 2: According to Corollary 1, if $C = 0$, then under the constraint $q_t \sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o + q_u \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) \leq P$ for the BC and the constraint $\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) \sigma^2 \leq P$ for the dual MAC, the two channels have the same SINR region. Since the achievable rates of user i in the MIMO-MAC and the MIMO-BC are $r_i^m = \sum_{j=1}^{N_r} \log(1 + \text{SINR}_{i,j}^m)$ and $r_i^b = \sum_{j=1}^{N_r} \log(1 + \text{SINR}_{i,j}^b)$, the rate regions of the two channels are the same. Therefore, Proposition 2 follows. ■

Note that due to the additional interference constraint, Problem 2 cannot be solved by using the established duality result in [7] and [8], in which only a single sum power constraint was considered. Our duality result in Proposition 2 can be thought as an extension of the duality results in [7] [8] to a multiple linear constraint case. Moreover, as will be shown in the following section, our duality result formulates a MIMO-MAC problem (Problem 4), which can be efficiently solved.

III. DUAL MAC WEIGHTED SUM RATE MAXIMIZATION PROBLEM

In this section, we propose an efficient algorithm to solve Problem 4. With the SIC scheme, the achievable rate of the k th user in the dual MIMO-MAC is given by

$$r_k^m = \log \frac{|\mathbf{R}_w + \sum_{j=1}^k \mathbf{H}_j \mathbf{Q}_j^m \mathbf{H}_j^\dagger|}{|\mathbf{R}_w + \sum_{j=1}^{k-1} \mathbf{H}_j \mathbf{Q}_j^m \mathbf{H}_j^\dagger|}. \quad (20)$$

For the MIMO-MAC, the *equally* weighted sum rate maximization is irrespective of the decoding order. However, in general the weighted sum rate maximization in the MIMO-MAC is affected by the decoding order. We thus need to consider the optimal decoding order of the SIC for the dual MIMO-MAC, and further need to consider the corresponding optimal encoding order of the DPC for the primal BC.

Let π be the optimal decoding order, which is a permutation on the SU index set $\{1, \dots, K\}$. It follows from [20] that the optimal user decoding order π for Problem 4 is the order such that $w_{\pi(1)} \geq w_{\pi(2)} \geq \dots \geq w_{\pi(K)}$ is satisfied. The following lemma presents the optimal decoding order of the SIC for the data streams within a SU (see Appendix B for the proof).

Lemma 1: The optimal data stream decoding order for a particular SU is arbitrary.

Due to the duality between the MIMO-BC and the MIMO-MAC, for Problem 3, the optimal encoding order for the DPC is the reverse of π . Because of the arbitrary encoding order for the data streams within a SU, if we choose a different encoding order for the BC, the MAC-to-BC mapping algorithm can give different results which yield the same objective value. Hence, the matrix \mathbf{Q}_i^b achieving the optimal objective value are not unique. With no loss of generality, we assume $w_1 \geq w_2 \geq \dots \geq w_K$ for notational convenience.

According to (20), the objective function of Problem 4 can be rewritten as

$$f(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m) := \sum_{i=1}^K \Delta_i \log |\mathbf{R}_w + \sum_{j=1}^i \mathbf{H}_j \mathbf{Q}_j^m \mathbf{H}_j^\dagger|, \quad (21)$$

where $\Delta_i := w_i - w_{i+1}$, and $w_{K+1} := 0$. Clearly, Problem 4 is a convex problem, which can be solved through standard convex optimization software packages directly. However, the standard convex optimization software does not exploit the special structure of the problem, and thus is computationally expensive. In the following, we develop a primal dual method based algorithm [18] to solve this problem.

We next rewrite Problem 4 as

$$\max_{\{\mathbf{Q}_i^m\}_{i=1}^K: \mathbf{Q}_i^m \succeq 0} f(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m) \quad \text{subject to} \quad \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) \leq P. \quad (22)$$

Recall that the positive semi-definiteness of \mathbf{Q}_i^m is equivalent to the positiveness of the eigenvalues of \mathbf{Q}_i^m , i.e., $q_{i,j} \geq 0$. Correspondingly, the Lagrangian function is

$$L(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m, \lambda, \delta_{i,j}) = f(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m) - \lambda \left(\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) - P \right) + \sum_{i=1}^K \sum_{j=1}^{M_i} \delta_{i,j} q_{i,j}, \quad (23)$$

where λ and $\delta_{i,j}$ are Lagrangian multipliers. According to the KKT conditions of (22), we have

$$\frac{\partial f(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m)}{\partial \mathbf{Q}_i^m} - \lambda \mathbf{I}_{N_r} + \sum_{j=1}^{M_i} \delta_{i,j} \mathbf{E}_{i,j} = 0, \quad (24)$$

$$\lambda \left(\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) - P \right) = 0, \quad (25)$$

$$\delta_{i,j} q_{i,j} = 0, \quad (26)$$

where $\mathbf{E}_{i,j} := \partial q_{i,j} / \partial \mathbf{Q}_i^m$. Notice that it is not necessary to compute the actual value of $\delta_{i,j}$ and $\mathbf{E}_{i,j}$, because if $\delta_{i,j} \neq 0$, then $q_{i,j} = 0$. Thus, the semi-definite constraint turns into $q_{i,j} = [q_{i,j}]^+$. Thus, we can assume $\delta_{i,j} = 0$.

The dual objective function of (22) is

$$g(\lambda) = \max_{\{\mathbf{Q}_i^m\}_{i=1}^K: \mathbf{Q}_i^m \geq 0} L(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m, \lambda). \quad (27)$$

Because the problem (22) is convex, it is equivalent to the following minimization problem

$$\min_{\lambda} g(\lambda) \quad \text{subject to} \quad \lambda \geq 0. \quad (28)$$

We outline the algorithm to solve the problem (28). We choose an initial λ and compute the value of $g(\lambda)$ (27), and then update λ according to the descent direction of $g(\lambda)$. The process repeats until the algorithm converges.

It is easy to observe that all the users share the same λ , and thus λ can be viewed as a water level in the water filling principle [21]. Once λ is fixed, the unique optimal set $\{\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m\}$ can be obtained through the gradient ascent algorithm. In each iterative step, \mathbf{Q}_i^m is updated sequentially according to its gradient direction of (23). Denote by $\mathbf{Q}_i^m(n)$ the matrix \mathbf{Q}_i^m at the n th iteration step. The gradient of each step is determined by

$$\nabla_{\mathbf{Q}_i^m}^{(n)} L := \frac{\partial f(\mathbf{Q}_1^m(n), \dots, \mathbf{Q}_{i-1}^m(n), \mathbf{Q}_i^m(n-1), \dots, \mathbf{Q}_K^m(n-1))}{\partial \mathbf{Q}_i^m(n-1)} - \lambda \mathbf{I}_{N_r}. \quad (29)$$

Thus, $\mathbf{Q}_i^m(n)$ can be updated according to

$$\mathbf{Q}_i^m(n) = \left[\mathbf{Q}_i^m(n-1) + t \nabla_{\mathbf{Q}_k^m}^{(n)} L \right]^+,$$

where t is the step size, and the notation $[\mathbf{A}]^+$ is defined as $[\mathbf{A}]^+ := \sum_j [\lambda_j]^+ \mathbf{v}_j \mathbf{v}_j^\dagger$ with λ_j and \mathbf{v}_j being the j th eigenvalue and the corresponding eigenvector of \mathbf{A} respectively. The gradient in (29) can be readily computed as

$$\frac{\partial f(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m)}{\partial \mathbf{Q}_k^m} = \sum_{j=k}^K \Delta_j \left(\mathbf{H}_k \mathbf{F}_j(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m)^{-1} \mathbf{H}_k^\dagger \right), \quad (30)$$

where $\mathbf{F}_j(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m) := \mathbf{R}_w + \sum_{i=1}^j \mathbf{H}_i^\dagger \mathbf{Q}_i^m \mathbf{H}_i$. We next need to determine the optimal λ . Since the Lagrangian function $g(\lambda)$ is convex over λ , the optimal λ can be obtained through the one-dimensional search. However, because $g(\lambda)$ is not necessarily differentiable, the gradient algorithm cannot be applied. Alternatively, the subgradient method can be used to find the optimal solution. In each iterative step, λ is updated according to the subgradient direction.

Lemma 2: The sub-gradient of $g(\lambda)$ is $P - \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m)$, where $\lambda \geq 0$, and \mathbf{Q}_i^m , $i = 1, \dots, K$, are the corresponding optimal covariance matrices for a fixed λ in (27).

Proof: The proof is provided in Appendix C. ■

Lemma 2 indicates that the value of λ should increase, if $\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) > P$, and vice versa. We now present our DIPA algorithm for solving Problem 4 in table I. The following proposition shows the convergence property of the DIPA algorithm.

Proposition 4: The DIPA algorithm converges to an optimal set of the MAC transmit signal covariance matrices.

Proof: The DIPA algorithm consists of the inner and outer loops. The inner loop is to compute \mathbf{Q}_i^m for $i = 1, \dots, K$. In each iterative step of the inner loop, we update \mathbf{Q}_i^m by fixing other \mathbf{Q}_j^m with $j \neq i$, and compute the corresponding gradient. The inner loop uses the gradient ascent algorithm, which converges to the optimal value due to its nondecreasing property and the convexity of the objective function. The outer loop is to compute the optimal Lagrangian multiplier λ in (28). Due to the convexity of the dual objective function [22], there is a unique λ achieving the optimal solution in (28). Hence, we can use an efficient one dimensional line bisection search [21], [12]. ■

Remark 2: In the previous work on the sum rate maximization [11], [12], [21], the covariance matrix of each user is the same as the single user water-filling covariance matrix in a point-to-point link with multiuser interference being treated as noise [23]. However, for the weighted sum rate maximization

problem, the optimal solution does not possess a water-filling structure. Thus, our DIPA algorithm does not obey the water-filling principle. In Section V, Example 1 compares the water-filling algorithm with the DIPA algorithm.

In the dual MIMO-MAC, according to (11), the transmit beamforming vectors $\mathbf{v}_{i,j}$ can be obtained by the eigenvalue decomposition. The corresponding receive beamforming vector at the BS, $\mathbf{u}_{i,j}$, is obtained by using the minimum mean square error (MMSE) algorithm. Throughout the proof of Proposition 3, we can see that when the same optimal solutions are achieved the primal BC and the dual MAC share the same beamforming vectors $\mathbf{u}_{i,j}$ and $\mathbf{v}_{i,j}$, and achieve the same SINR region, i.e., $\text{SINR}_{i,j}^{\text{b}} = \text{SINR}_{i,j}^{\text{m}}$. Based on this observation, we can compute the power allocated to the BC beamforming direction $\mathbf{u}_{i,j}$, $p_{i,j}$, and obtain the signal covariance matrix, $\mathbf{Q}_i^{\text{b}} = \sum_{j=1}^{N_r} p_{i,j} \mathbf{u}_{i,j} \mathbf{u}_{i,j}^{\dagger}$. The MAC-to-BC covariance matrix mapping allows us to obtain the optimal BC covariance matrices for Problem 3 by solving Problem 4.

IV. A COMPLETE SOLUTION TO THE CR MIMO-BC WEIGHTED SUM RATE PROBLEM

We are now ready to present a complete algorithm to solve Problem 2. The Lagrangian dual objective function of Problem 2 can be rewritten as follows

$$g(q_t, q_u) = \max_{\{\mathbf{Q}_i^{\text{b}}\}_{i=1}^K: \mathbf{Q}_i^{\text{b}} \succeq 0} \sum_{i=1}^K w_i r_i^{\text{b}}, \quad (31)$$

where the maximization is subject to the constraint $q_t (\sum_{i=1}^K \mathbf{h}_o^{\dagger} \mathbf{Q}_i^{\text{b}} \mathbf{h}_o - P_t) + q_u (\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^{\text{b}}) - P_u) \leq 0$. Problem 2 is equivalent to the following problem

$$\min_{q_t, q_u} g(q_t, q_u), \quad \text{subject to } q_t \geq 0 \text{ and } q_u \geq 0.$$

Applying the BC-MAC duality in Section II-B and the DIPA algorithm in Section III, $g(q_t, q_u)$ can be obtained. The remaining task is to determine the optimal q_t and q_u . Since $g(q_t, q_u)$ is not necessarily differentiable, we search the optimal q_t and q_u through the subgradient algorithm; that is, in each iterative step, we update the vector $[q_t, q_u]$ according to the subgradient direction $\mathbf{s} = [s_1, s_2]$ of $g(q_t, q_u)$.

Lemma 3: The subgradient of $g(q_t, q_u)$ is $[P_t - \sum_{i=1}^K \mathbf{h}_o^{\dagger} \mathbf{Q}_i^{\text{b}} \mathbf{h}_o, P_u - \sum_{i=1}^K \text{tr}(\mathbf{Q}_i^{\text{b}})]$, where $q_t \geq 0, q_u \geq 0$, and $\mathbf{Q}_i^{\text{b}}, i = 1, \dots, K$, are the corresponding optimal covariance matrices for the problem (31).

The proof of Lemma 3 is similar to that of Lemma 2, and is omitted here. It has been shown in [24] that with a constant step size, the subgradient algorithm converges to a value that is within a small range of the optimal value, i.e.,

$$\lim_{n \rightarrow \infty} |q_t^{(n)} - q_t^*| < \epsilon, \quad \text{and}, \quad \lim_{n \rightarrow \infty} |q_u^{(n)} - q_u^*| < \epsilon, \quad (32)$$

where q_t^* and q_u^* denote the optimal values, and $q_t^{(n)}$ and $q_u^{(n)}$ denote the values of q_t and q_u at the n th step of the subgradient algorithm, respectively. This implies that the subgradient method finds an ϵ -suboptimal point within a finite number of steps. The number ϵ is a decreasing function of the step size.

We next describe the algorithm to solve Problem 2 in table II. As a summary, the flow chart of the SIPA algorithm is depicted in Fig. 3. We shows that the SIPA algorithm converges to the optimal solution of Problem 1 in the following proposition.

Proposition 5: The SIPA algorithm converges to the globally optimal solution of Problem 1.

Proof: The Lagrangian function of Problem 1 is given by

$$L(\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b, \lambda_1, \lambda_2) = \sum_{i=1}^K w_i r_i^b - \lambda_1 \left(\sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o - P_t \right) - \lambda_2 \left(\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P_u \right), \quad (33)$$

and the Lagrangian function of Problem 2 is given by

$$L_1(\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b, \lambda, q_t, q_u) = \sum_{i=1}^K w_i r_i^b - \lambda \left(q_t \left(\sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o - P_t \right) - q_u \left(\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P_u \right) \right). \quad (34)$$

Let \bar{q}_t , \bar{q}_u , $\bar{\lambda}$, and $\bar{\mathbf{Q}}_i$ be the optimal values of $L_1(\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b, \lambda, q_t, q_u)$, when the algorithm converges. We thus have

$$\left. \frac{\partial L_1(\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b, \lambda, q_t, q_u)}{\partial \mathbf{Q}_i^b} \right|_{\{\bar{\mathbf{Q}}_i\}_{i=1}^K, \bar{\lambda}, \bar{q}_t, \bar{q}_u} = 0,$$

$|\bar{q}_t (\sum_{i=1}^K \mathbf{h}_o^\dagger \bar{\mathbf{Q}}_i \mathbf{h}_o - P_t)| = 0$, and $|\bar{q}_u (\sum_{i=1}^K \text{tr}(\bar{\mathbf{Q}}_i) - P_u)| = 0$. This means that $\bar{\mathbf{Q}}_i$ is a locally optimal solution.

According to (33), if we select $\tilde{\lambda}_1 = \bar{\lambda} \bar{q}_t$, $\tilde{\lambda}_2 = \bar{\lambda} \bar{q}_u$, and $\tilde{\mathbf{Q}}_i = \bar{\mathbf{Q}}_i$, then $\tilde{\lambda}_1$, $\tilde{\lambda}_2$, and $\tilde{\mathbf{Q}}_i$ satisfy the KKT conditions of Problem 1 and thus are the locally optimal variables. Suppose that there exists an optimal set of $\hat{\lambda}_1$, $\hat{\lambda}_2$, and $\hat{\mathbf{Q}}_i$ such that $L(\hat{\mathbf{Q}}_1, \dots, \hat{\mathbf{Q}}_K, \hat{\lambda}_1, \hat{\lambda}_2) > L(\tilde{\mathbf{Q}}_1, \dots, \tilde{\mathbf{Q}}_K, \tilde{\lambda}_1, \tilde{\lambda}_2)$. Clearly, this optimal set of $\hat{\lambda}_1$, $\hat{\lambda}_2$, and $\hat{\mathbf{Q}}_i$ satisfy the KKT conditions of Problem 1. In the sequel, we will derive a contradiction.

First, we can write

$$L(\tilde{\mathbf{Q}}_1, \dots, \tilde{\mathbf{Q}}_K, \tilde{\lambda}_1, \tilde{\lambda}_2) \geq L(\hat{\mathbf{Q}}_1, \dots, \hat{\mathbf{Q}}_K, \tilde{\lambda}_1, \tilde{\lambda}_2). \quad (35)$$

Suppose that the inequality (35) does not hold, i.e., $L(\tilde{\mathbf{Q}}_1, \dots, \tilde{\mathbf{Q}}_K, \tilde{\lambda}_1, \tilde{\lambda}_2) < L(\hat{\mathbf{Q}}_1, \dots, \hat{\mathbf{Q}}_K, \tilde{\lambda}_1, \tilde{\lambda}_2)$. Then, according to the BC-MAC duality in Section II-B, an objective value of (7) which is larger than $L(\tilde{\mathbf{Q}}_1, \dots, \tilde{\mathbf{Q}}_K, \tilde{\lambda}_1, \tilde{\lambda}_2)$, can be found for the fixed \bar{q}_t and \bar{q}_u . However, from Proposition 4, the DIPA algorithm converges the optimal solution. It is a contradiction.

Secondly, according to the KKT conditions of Problem 1, we have

$$\hat{\lambda}_1 \left(\sum_{i=1}^K \mathbf{h}_o^\dagger \hat{\mathbf{Q}}_i^b \mathbf{h}_o - P_t \right) = 0, \quad \hat{\lambda}_2 \left(\sum_{i=1}^K \text{tr}(\hat{\mathbf{Q}}_i^b) - P_u \right) = 0. \quad (36)$$

We thus can write:

$$L(\hat{\mathbf{Q}}_1, \dots, \hat{\mathbf{Q}}_K, \tilde{\lambda}_1, \tilde{\lambda}_2) \geq L(\hat{\mathbf{Q}}_1, \dots, \hat{\mathbf{Q}}_K, \hat{\lambda}_1, \hat{\lambda}_2). \quad (37)$$

Combining (37) and (35), we have

$$L(\tilde{\mathbf{Q}}_1, \dots, \tilde{\mathbf{Q}}_K, \tilde{\lambda}_1, \tilde{\lambda}_2) \geq L(\hat{\mathbf{Q}}_1, \dots, \hat{\mathbf{Q}}_K, \hat{\lambda}_1, \hat{\lambda}_2). \quad (38)$$

This contradicts with our previous assumption. ■

Remark 3: The algorithm can be extended to the multiple PU case in the following manner. Assume that there are N PUs. Problem 2 becomes

$$\begin{aligned} & \min_{q_{t,j} \geq 0, q_u \geq 0} \max_{\{\mathbf{Q}_i^b\}_{i=1}^K: \mathbf{Q}_i^b \succeq 0} \sum_{i=1}^K w_i r_i^b, \\ & \text{subject to } \sum_{j=1}^N q_{t,j} \left(\sum_{i=1}^K \mathbf{h}_{o,j}^\dagger \mathbf{Q}_i^b \mathbf{h}_{o,j} - P_{t,j} \right) + q_u \left(\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P_u \right) \leq 0, \end{aligned} \quad (39)$$

where $q_{t,j}$ is the auxiliary variable for the j th PU, $\mathbf{h}_{o,j}$ is the channel response from the BS to the j th PU, and $P_{t,j}$ is the interference threshold of the j th PU. The role of auxiliary variables $q_{t,j}$ is similar to that of q_t in the single PU case. It is thus straightforward to modify the SIPA algorithm to solve the problem for the multiple PU case. Moreover, it should be noted that the multiple interference constraints of the problem (39) can be transformed to the per-antenna power constraints [17] by setting $\mathbf{h}_{o,j}$, $j = 1, \dots, N_t$, to be the j th column of the identity matrix. Not limited by the sum rate maximization problem with interference power constraints, the method proposed in this paper can be easily applied to solve the transmitter optimization problem (e.g. beamforming optimization or power minimization) for MIMO BC with multiple linear power constraints.

V. SIMULATION RESULTS

In this section, we provide the simulation results to show the effectiveness of the proposed algorithm. In the simulations, for simplicity, we assume that the BS is at the same distance, l_1 , to all SUs, and the same distance, $l_2^{(n)}$, to PU_n . In the single PU case, we will drop the superscript and simply use notation l_2 . Suppose that the same path loss model can be used to describe the transmissions from the BS to the SUs and to the PUs, and the pass loss exponent is 4. The elements of matrix \mathbf{H} are assumed to be circularly

symmetric complex Gaussian (CSCG) RVs with mean zero and variance 1, and \mathbf{h}_o can be modeled as $\mathbf{h}_o = (l_1/l_2)^2 \mathbf{a}_n$, where \mathbf{a}_n is a $N_t \times 1$ vector whose elements are CSCG RVs with mean zero and variance 1. The noise covariance matrix at the BS is assumed to be the identity matrix, and the sum power and interference power are defined in dB relative to the noise power, and P_t is chosen to be 0 dB. For all cases, we choose $l_1 = l_2$, except for explicitly stated.

Example 1: In Fig. 4, we examine the validity of the DIPA algorithm. In this example, we choose $K = 1$ (a single SU case), $N_t = 4$, $N_r = 4$, and $P_u = 10$ dB. It is well known that the optimal transmit signal covariance matrix can be obtained through the water-filling principle [23]. As can be observed from Fig. 4, in several iterations, the DIPA algorithm converges to the optimal solution obtained by using the water-filling principle.

Example 2: In Fig. 5, we show the convergence property of the DIPA algorithm. In this example, we choose $K = 20$ and $P_u = 10$ dB. It can be observed from this figure that the algorithm converges to the optimal solution within several iteration steps.

Example 3: In Figs. 6 and 7, we consider a SU MIMO-BC network with $K = 5$, $N_t = 5$, $N_r = 3$, and $P_u = 13$ dB. In this example, the SUs with $w_1 = 5$ and $w_i = 1$, $i = 2, \dots, K$ are assumed to share the same spectrum band with two PUs. Fig. 6 plots the weighted sum rate versus the number of iterations of the SIPA algorithm for step sizes $t = 0.1$ and $t = 0.01$. As can be seen from the figure, the step size affects the accuracy and convergence speed of the algorithm. Fig. 7 plots the sum power at the BS and the interference power at the PUs versus the number of iterations. It can be seen from the figure that the sum power and the interference power approach to $P_u = 13$ dB and $P_t = 0$ dB respectively when the SIPA algorithm converges. This implies that the sum power and interference constraints are satisfied with equalities when the SIPA algorithm converges.

Example 4: Fig. 8 plots the achievable sum rates versus the sum power in the single PU case and the case with no PU. We choose $K = 5$, $N_t = 5$, and $N_r = 3$. As can be seen from Fig. 8, in the low sum power regime, the achievable sum rate in the case with no PU is quite close to the one in the single PU case while in the high sum power regime, the achievable sum rate in the case with no PU is much higher than the one in the single PU case. This is because the additional constraint reduces the degrees of freedom of the system.

VI. CONCLUSIONS

In this paper, we developed a new BC-MAC duality result, which can be viewed as an extension of existing duality results developed under either a sum power constraint or per-antenna power constraints. Exploiting this duality result, we proposed an efficient algorithm to solve the CR MIMO-BC weighted sum rate maximization problem. We further showed that the proposed algorithm converges to the globally optimal solution.

APPENDIX

A. Proof of Proposition 1: The proof consists of two parts. In the first part, we show that either optimal solution is feasible for both problems. In the second part, we show that Problem 1 and Problem 2 have the same solution.

The Lagrangian function of Problem 1 is

$$L_1(\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b, \lambda_t, \lambda_u) = \sum_{i=1}^K w_i r_i^b - \lambda_t \left(\sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o - P_t \right) - \lambda_u \left(\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P_u \right), \quad (40)$$

where λ_t and λ_u are the Lagrangian multipliers. The optimal objective value is

$$\min_{\lambda_t, \lambda_u} \max_{\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b} L_1(\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b, \lambda_t, \lambda_u). \quad (41)$$

Assume the optimal variables are $\bar{\lambda}_t$, $\bar{\lambda}_u$ and $\bar{\mathbf{Q}}_1^b, \dots, \bar{\mathbf{Q}}_K^b$, and the corresponding optimal value is \bar{C} .

The Lagrangian function of Problem 2 is:

$$L_2(\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b, q_t, q_u, \lambda) = \sum_i w_i r_i^b - \lambda \left(q_t \left(\sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o - P_t \right) + q_u \left(\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P_u \right) \right), \quad (42)$$

where λ is the Lagrangian multiplier. The optimal objective value is

$$\min_{q_t, q_u, \lambda} \max_{\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b} L_2(\mathbf{Q}_1^b, \dots, \mathbf{Q}_K^b, q_t, q_u, \lambda). \quad (43)$$

Suppose that the optimal variables are \tilde{q}_t , \tilde{q}_u , $\tilde{\lambda}$, and $\tilde{\mathbf{Q}}_i^b$, $i = 1, \dots, K$, and the corresponding optimal objective value is \tilde{C} . We just need to prove $\bar{C} = \tilde{C}$.

We now present the first part of the proof. According to the KKT condition of Problem 2, we have

$$\frac{\partial L_2(\tilde{\lambda}, \tilde{\mathbf{Q}}_1^b, \dots, \tilde{\mathbf{Q}}_K^b, \tilde{q}_t, \tilde{q}_u)}{\partial q_t} = \tilde{\lambda} \left(\sum_{i=1}^K \mathbf{h}_o^\dagger \tilde{\mathbf{Q}}_i^b \mathbf{h}_o - P_t \right) = 0, \quad (44)$$

$$\frac{\partial L_2(\tilde{\lambda}, \tilde{\mathbf{Q}}_1^b, \dots, \tilde{\mathbf{Q}}_K^b, \tilde{q}_t, \tilde{q}_u)}{\partial q_u} = \tilde{\lambda} \left(\sum_{i=1}^K \text{tr}(\tilde{\mathbf{Q}}_i^b) - P_u \right) = 0. \quad (45)$$

Recall that the Lagrangian multiplier $\tilde{\lambda}$ is non-negative. Furthermore, if $\tilde{\lambda} = 0$, we have $\tilde{q}_t(\sum_{i=1}^K \mathbf{h}_o^\dagger \tilde{\mathbf{Q}}_i \mathbf{h}_o - P_t) + \tilde{q}_u(\sum_{i=1}^K \text{tr}(\tilde{\mathbf{Q}}_i) - P_u) < 0$ from the KKT conditions. Under this case, it is easy to prove by contradiction that $\tilde{\mathbf{Q}}_i^b$ is not the optimal solution. Thus, we always have $\tilde{\lambda} > 0$ and can readily conclude that $\sum_{i=1}^K \mathbf{h}_o^\dagger \tilde{\mathbf{Q}}_i^b \mathbf{h}_o = P_t$ and $\sum_{i=1}^K \text{tr}(\tilde{\mathbf{Q}}_i^b) = P_u$ are satisfied simultaneously. The optimal solution of Problem 2 is also a feasible solution of Problem 1. On the other hand, it is obvious that the feasible solution for Problem 1 is also the feasible solution for Problem 2.

We next prove the second part by using contradiction. Let us first suppose $\bar{C} > \tilde{C}$. For (42), if we select $\mathbf{Q}_i^b = \tilde{\mathbf{Q}}_i^b$ for $i = 1, \dots, K$, $\lambda = 1$, $q_t = \bar{\lambda}_t$ and $q_u = \bar{\lambda}_u$, then $L_2 = \bar{C} > \tilde{C}$. It contradicts to the fact that \tilde{C} is the optimal objective value for (43).

We now assume $\bar{C} < \tilde{C}$. Recall that $\tilde{\lambda} \neq 0$, for (42). If we select $\mathbf{Q}_i^b = \tilde{\mathbf{Q}}_i^b$ for $i = 1, \dots, K$, $\lambda_t = \tilde{\lambda} \tilde{q}_t$ and $\lambda_u = \tilde{\lambda} \tilde{q}_u$, then $L_1 = \tilde{C} > \bar{C}$, which contradicts with the fact that \bar{C} is the optimal objective value for (41). Therefore, the optimal solutions for Problem 2 and Problem 1 are the same. ■

B. Proof of Lemma 1: According to previous discussions, the signal from each SU is divided into several data streams. We now show that the optimal encoding order of these data streams are arbitrary. It is well known that the optimal objective value of the MAC equally weighted sum rate problem can be achieved by adopting any ordering [21] [11] [12]; that is, when all the users have the same weights, the optimal solution of the weighted sum rate maximization problem is independent of the decoding order. Analogously, the data streams within a SU share the same weight. Thus, an arbitrary encoding order of those data streams within a SU can achieve the optimal solution. ■

C. Proof of Lemma 2: Let s be the sub-gradient of $g(\tilde{\lambda})$. For a given $\check{\lambda} \geq 0$, the subgradient s of $g(\tilde{\lambda})$ satisfies $g(\check{\lambda}) \geq g(\tilde{\lambda}) + s(\check{\lambda} - \tilde{\lambda})$, where $\check{\lambda}$ is any feasible value. Let $\check{\mathbf{Q}}_i^m$, $i = 1, \dots, K$, be the optimal covariance matrices in (27) for $\lambda = \check{\lambda}$, and $\tilde{\mathbf{Q}}_i^m$, $i = 1, \dots, K$, be the optimal covariance matrices in (27) for $\lambda = \tilde{\lambda}$. We express $g(\check{\lambda})$ as

$$\begin{aligned} g(\check{\lambda}) &= \max_{\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m} \left(f(\mathbf{Q}_1^m, \dots, \mathbf{Q}_K^m) - \check{\lambda} \left(\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) - P \right) \right) \\ &= f(\check{\mathbf{Q}}_1^m, \dots, \check{\mathbf{Q}}_K^m) - \check{\lambda} \left(\sum_{i=1}^K \text{tr}(\check{\mathbf{Q}}_i^m) - P \right) \\ &\geq f(\tilde{\mathbf{Q}}_1^m, \dots, \tilde{\mathbf{Q}}_K^m) - \check{\lambda} \left(\sum_{i=1}^K \text{tr}(\tilde{\mathbf{Q}}_i^m) - P \right) \end{aligned}$$

$$\begin{aligned}
&= f(\tilde{\mathbf{Q}}_1^m, \dots, \tilde{\mathbf{Q}}_K^m) - \tilde{\lambda} \left(\sum_{i=1}^K \text{tr}(\tilde{\mathbf{Q}}_i^m) - P \right) + \tilde{\lambda} \left(\sum_{i=1}^K \text{tr}(\tilde{\mathbf{Q}}_i^m) - P \right) - \check{\lambda} \left(\sum_{i=1}^K \text{tr}(\tilde{\mathbf{Q}}_i^m) - P \right) \\
&= g(\tilde{\lambda}) + \left(P - \sum_{i=1}^K \text{tr}(\tilde{\mathbf{Q}}_i^m) \right) (\check{\lambda} - \tilde{\lambda}),
\end{aligned}$$

where $s := P - \sum_{i=1}^K \text{tr}(\tilde{\mathbf{Q}}_i^m)$ is the subgradient of $g(\tilde{\lambda})$. This concludes the proof. \blacksquare

REFERENCES

- [1] J. Mitola and G. Q. Maguire, "Cognitive radios: Making software radios more personal," *IEEE Personal Communications*, vol. 6, no. 4, pp. 13–18, Aug. 1999.
- [2] S. Haykin, "Cognitive radio: Brain-empowered wireless communications," *IEEE J. Select. Areas Commun.*, vol. 23, no. 2, pp. 201–202, Feb. 2005.
- [3] M. Gastpar, "On capacity under receive and spatial spectrum-sharing constraints," *IEEE Trans. Inform. Theory*, vol. 53, no. 2, pp. 471–487, Feb. 2007.
- [4] A. Ghasemi and E. S. Sousa, "Fundamental limits of spectrum-sharing in fading environments," *IEEE Trans. Wireless Commun.*, vol. 6, no. 2, pp. 649–658, Feb. 2007.
- [5] Y.-C. Liang, Y. Zeng, E. Peh, and A. Hoang, "Sensing-throughput tradeoff for cognitive radio networks," *IEEE Trans. Wireless Commun.*, vol. 7, no. 4, pp. 1326–1337, Apr. 2008.
- [6] Z. Quan, S. Cui, and A. Sayed, "Optimal linear cooperation for spectrum sensing in cognitive radio networks," *IEEE J. Select. Topics in Signal Processing*, vol. 2, no. 1, pp. 28–40, Feb. 2008.
- [7] P. Viswanath and D. N. C. Tse, "Sum capacity of the vector Gaussian broadcast channel and uplink-downlink duality," *IEEE Trans. Inform. Theory*, vol. 49, no. 8, pp. 1912–1921, Aug. 2003.
- [8] S. Vishwanath, N. Jindal, and A. Goldsmith, "Duality, achievable rates, and sum-rate capacity of Gaussian MIMO broadcast channels," *IEEE Trans. Inform. Theory*, vol. 49, no. 10, pp. 2658–2668, Oct. 2003.
- [9] W. Yu and J. Cioffi, "Sum capacity of Gaussian vector broadcast channels," *IEEE Transactions on information theory*, vol. 50, pp. 1875–1892, Sept. 2004.
- [10] H. Weingarten, Y. Steinberg, and S. Shamai, "The capacity region of the Gaussian multiple-input multiple-output broadcast channel," *IEEE Trans. Inform. Theory*, vol. 52, no. 9, pp. 3936–64, Sept. 2006.
- [11] N. Jindal, W. Rhee, S. Vishwanath, S. A. Jafar, and A. Goldsmith, "Sum power iterative water-filling for multi-antenna Gaussian broadcast channels," *IEEE Trans. Inform. Theory*, vol. 51, no. 4, pp. 1570–1580, Apr. 2005.
- [12] W. Yu, "Sum-capacity computation for the Gaussian vector broadcast channel via dual decomposition," *IEEE Trans. Inform. Theory*, vol. 52, no. 2, pp. 754–759, Feb. 2006.
- [13] M. Kobayashi and G. Caire, "An iterative water-filling algorithm for maximum weighted sum-rate of Gaussian MIMO-BC," *IEEE J. Select. Areas Commun.*, vol. 24, no. 8, pp. 1640–1646, Aug. 2006.
- [14] L. Jia and T. Hou, "Maximum weighted sum rate of multi-antenna broadcast channels," 2007. [Online]. Available: <http://arXiv:cs/0703111v1>.
- [15] F. Rashid-Farrokhi, L. Tassiulas, and K. Liu, "Joint optimal power control and beamforming in wireless networks using antenna arrays," *IEEE Trans. Commun.*, vol. 46, no. 10, pp. 1313–1324, Oct. 1998.
- [16] W. Yu, "Uplink-downlink duality via minimax duality," *IEEE Trans. Inform. Theory*, vol. 52, no. 2, pp. 361–374, Feb. 2006.
- [17] W. Yu and T. Lan, "Transmitter optimization for the multi-antenna downlink with per-antenna power constraints," *IEEE Trans. Signal Processing*, vol. 55, no. 6, pp. 2646–2660, June 2007.
- [18] D. G. Luenberger, *Optimization by vector space methods*. New York: John Wiley, 1969.
- [19] A. Wiesel, Y. C. Eldar, and S. Shamai, "Linear precoding via conic optimization for fixed MIMO receivers," *IEEE Trans. Signal Processing*, vol. 54, no. 1, pp. 161–176, Jan. 2006.
- [20] D. Tse and S. Hanly, "Multiaccess fading channels-part I: Polymatroid structure, optimal resource allocation and throughput capacities," *IEEE Trans. Inform. Theory*, vol. 44, no. 7, pp. 2796–2815, Nov. 1998.
- [21] W. Yu, W. Rhee, S. Boyd, and J. M. Cioffi, "Iterative water-filling for Gaussian vector multiple-access channels," *IEEE Trans. Inform. Theory*, vol. 50, no. 1, pp. 145–152, Jan. 2004.
- [22] S. Boyd and L. Vandenberghe, *Convex Optimization*. Cambridge, UK: Cambridge University Press, 2004.
- [23] I. E. Telatar, "Capacity of multi-antenna Gaussian channels," *European Trans. on Telecomm.*, vol. 10, no. 6, pp. 585–595, Oct. 1999.
- [24] S. Boyd, L. Xiao, and A. Mutapcic, "Subgradient methods," 2003. [Online]. Available: http://mit.edu/6.976/www/notes/subgrad_method.pdf.

TABLE I
DECOUPLED ITERATIVE POWER ALLOCATION (DIPA) ALGORITHM.

```

Initialize  $\lambda_{\min}$  and  $\lambda_{\max}$ ;
repeat
   $\lambda = (\lambda_{\min} + \lambda_{\max})/2$ 
  repeat, initialize  $\mathbf{Q}_1^m(0), \dots, \mathbf{Q}_K^m(0)$ ,  $n = 1$ 
    for  $i = 1, \dots, K$ ,
       $\mathbf{Q}_i^m(n) = [\mathbf{Q}_i^m(n-1) + t\nabla_{\mathbf{Q}_i^m}^{(n)} L]^+$ ,
    end for,
     $n \leftarrow n + 1$ ,
  until  $\mathbf{Q}_k^m$  for  $k = 1, \dots, K$  converge, i.e.,  $\|\nabla_{\mathbf{Q}_i^m}^{(n)} L\|^2 \leq \hat{\epsilon}$  for a small preset  $\hat{\epsilon}$ .
  if  $\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) > P$ , then  $\lambda_{\min} = \lambda$ , elseif  $\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^m) < P$ , then  $\lambda_{\max} = \lambda$ ;
until  $|\lambda_{\min} - \lambda_{\max}| \leq \epsilon$ 

```

TABLE II
SUBGRADIENT ITERATIVE POWER ALLOCATION (SIPA) ALGORITHM.

```

Initialization:  $q_t^{(1)}, q_u^{(1)}$ ,  $n = 1$ ;
repeat
  Find the optimal solution of the dual MAC Problem 4 through the DIPA algorithm;
  Find the solution of the BC problem (31) through the MAC-to-BC mapping;
  Update  $q_t^{(n)}$  and  $q_u^{(n)}$  through a subgradient algorithm  $q_t^{(n+1)} = q_t^{(n)} + t(\sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o - P_t)$ ,
    and  $q_u^{(n+1)} = q_u^{(n)} + t(\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P_u)$ , where  $t$  denotes the step size of the subgradient algorithm;
   $n \leftarrow n + 1$ ;
Stop when  $|q_t^{(n)}(\sum_{i=1}^K \mathbf{h}_o^\dagger \mathbf{Q}_i^b \mathbf{h}_o - P_t)| \leq \epsilon$  and  $|q_u^{(n)}(\sum_{i=1}^K \text{tr}(\mathbf{Q}_i^b) - P_u)| \leq \epsilon$  are satisfied simultaneously,

```

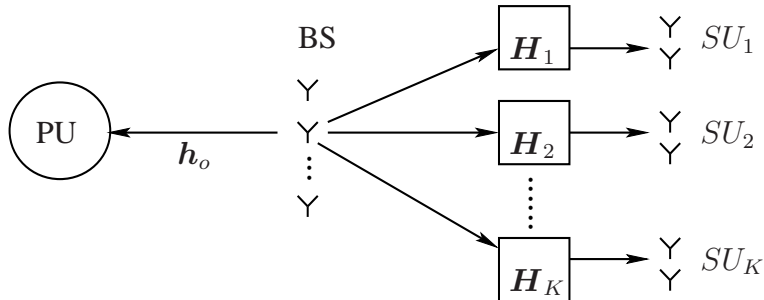


Fig. 1. The system model for MIMO-BC based cognitive radio networks. There are K SUs. The BS of the SUs has N_t transmit antennas, and each SU is equipped with N_r receive antennas.

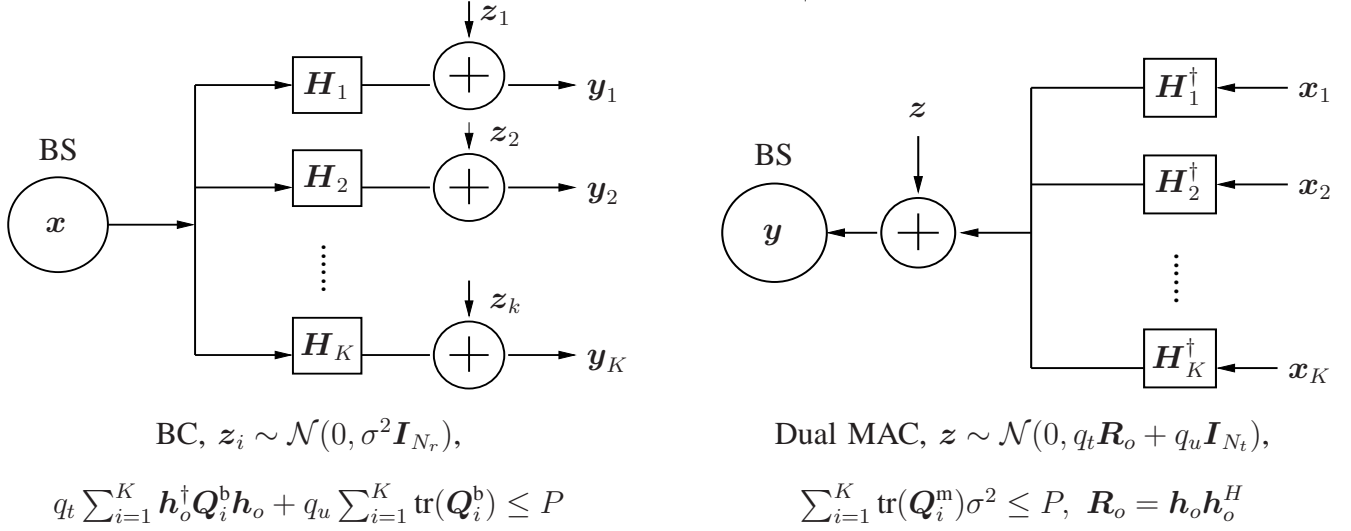


Fig. 2. The system models for Problem 3 and Problem 4, where q_t and q_u are constant, and $\mathbf{R}_o = \mathbf{h}_o \mathbf{h}_o^\dagger$.

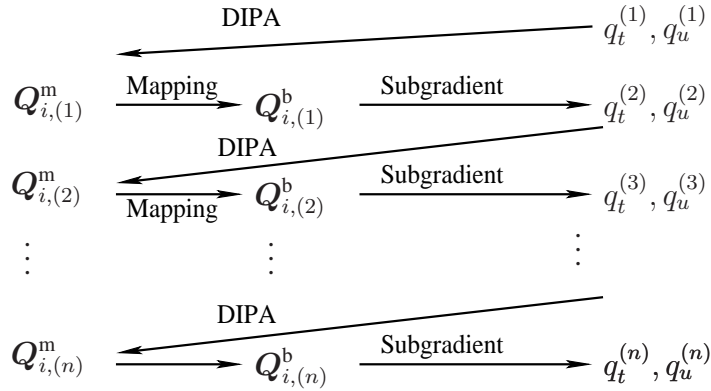


Fig. 3. The flow chart for the SIPA algorithm, where $Q_{i,(n)}^b$ and $Q_{i,(n)}^m$ denote the transmit signal covariance matrices of SU_i for the BC and MAC at the n th step, respectively.

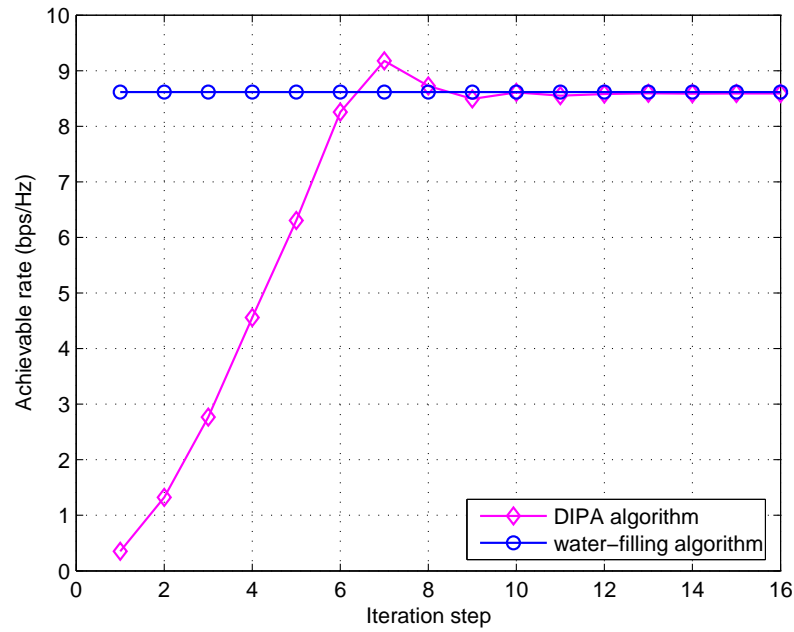


Fig. 4. Comparison of the optimal achievable rates obtained by the DIPA and the water-filling algorithm in a MIMO channel ($N_t = N_r = 4$, $K = 1$ and $P_u = 10$ dB).

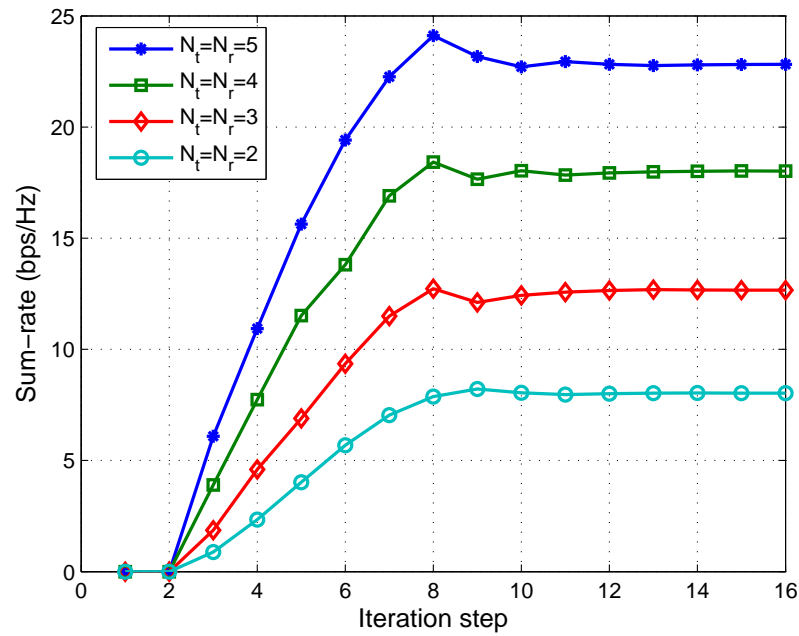


Fig. 5. Convergence behavior of the DIPA algorithm ($K = 20$ and $P_u = 10$ dB).

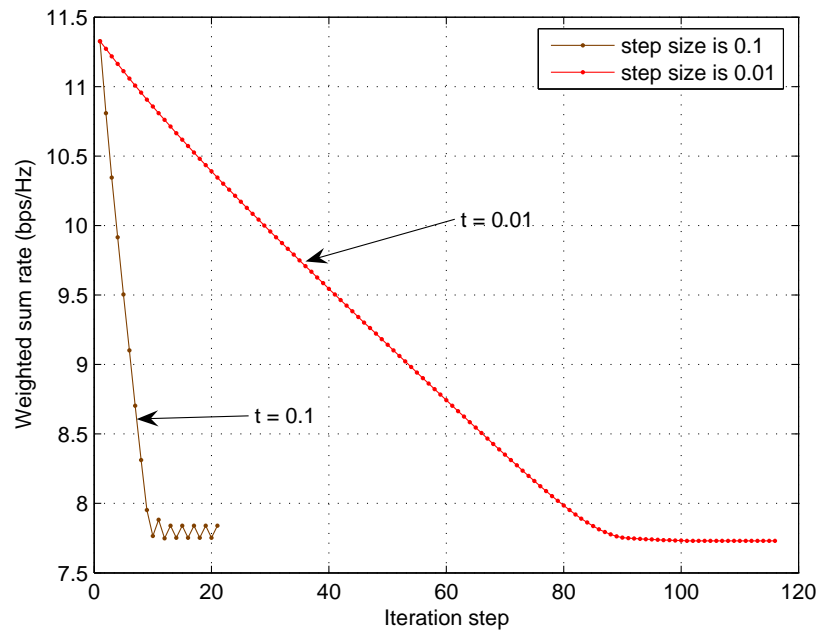


Fig. 6. Convergence behavior of the SIPA algorithm ($N_t = 5$, $K = 5$, $N_r = 3$, $w_1 = 5$, and $w_i = 1$, for $i \neq 1$).

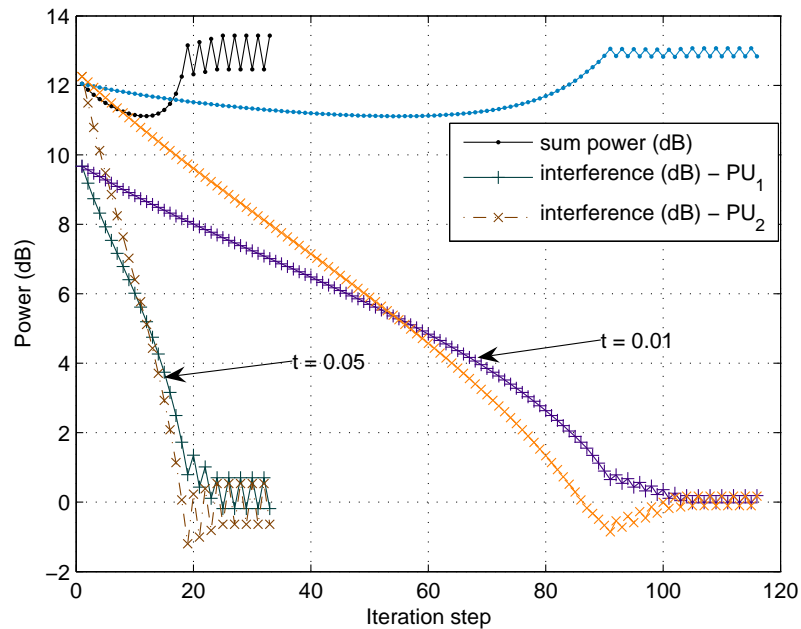


Fig. 7. The convergence behavior of the sum power at the BS and the interference at the PU for the SIPA algorithm ($N_t = 5$, $K = 5$, $N_r = 3$, $w_1 = 5$, and $w_i = 1$ with $i \neq 1$).

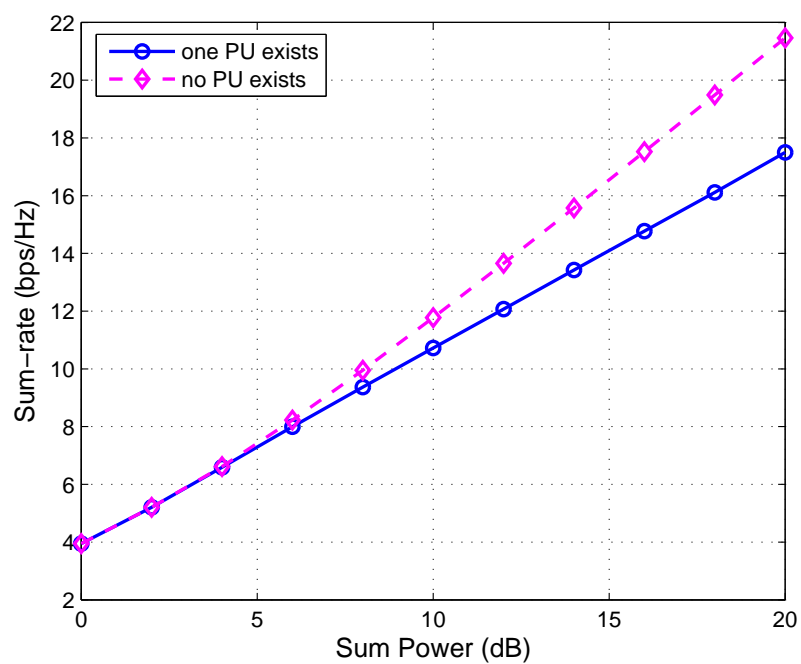


Fig. 8. Achievable sum rates versus sum power in the single PU case and the case with no PU ($N_t = 5$, $K = 5$, $N_r = 3$).